



Massachusetts  
Institute of  
Technology



# Approximation theory of graph neural networks

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MIT, TU Munich

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# Roadmap and summary

1. Approximation theory
  - a. Neural network approximation
  - b. Challenges to graph input
2. Graph embeddings and graph limits
  - a. Graphons and the limit of dense graph sequences
  - b. Challenges to sparse graph sequences

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# Deep learning framework



Ground truth  $f^* : \mathcal{X} \rightarrow \mathcal{Y}$  that is:

- Continuous/smooth
- Regular

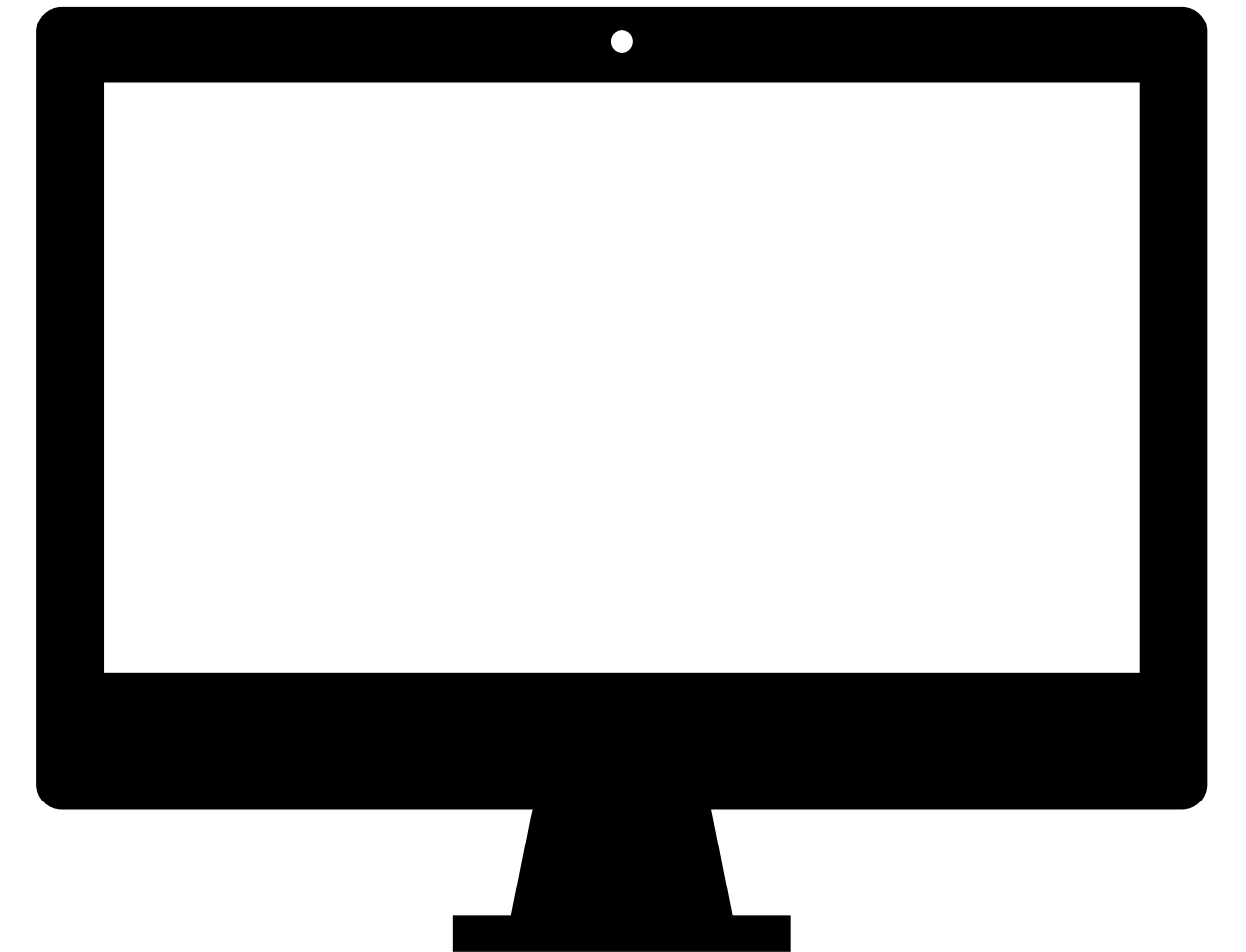
...

Model

$f_\theta : \mathcal{X} \rightarrow \mathcal{Y}$  that is:

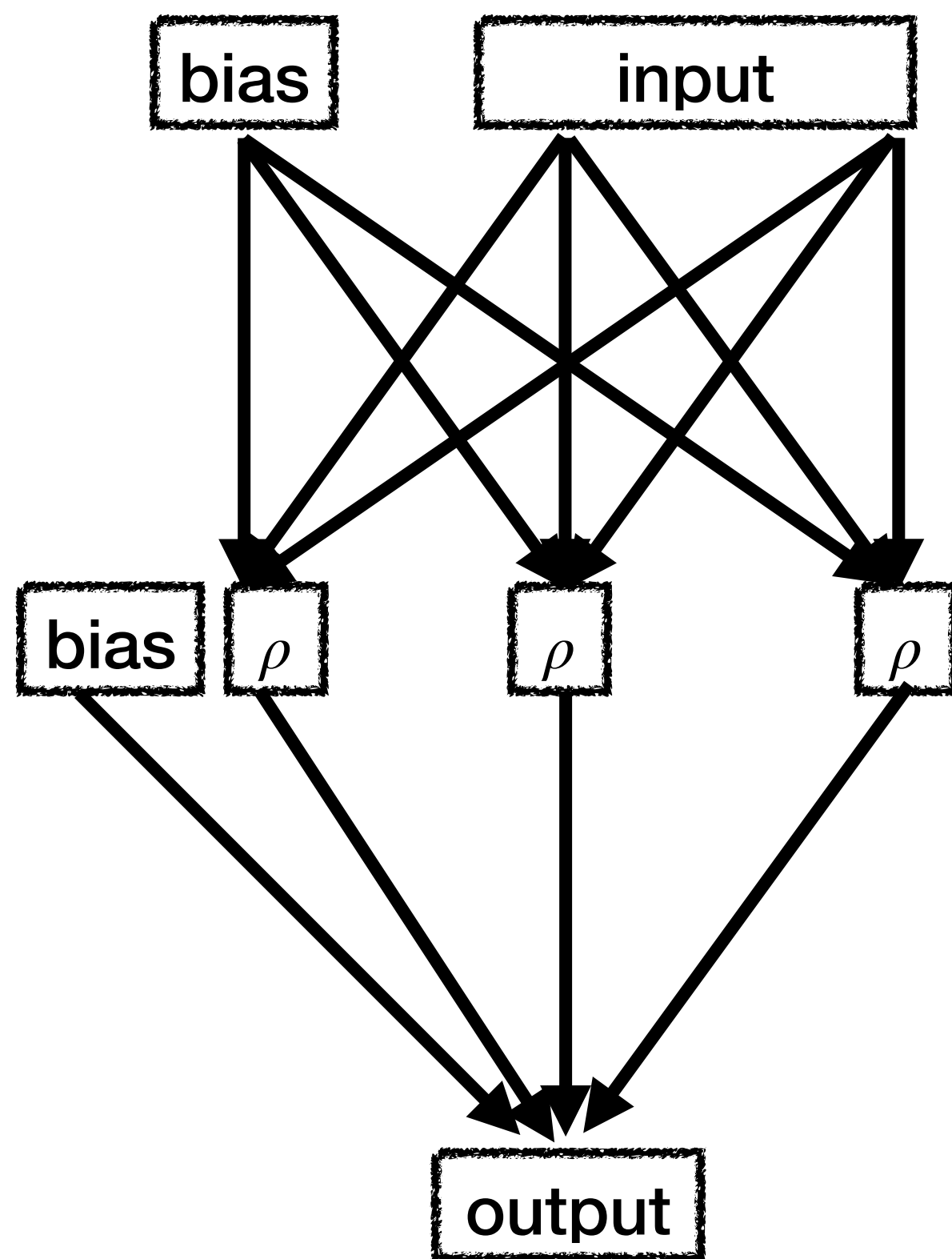
- Universal
- Optimization/GPU-friendly
- Efficient

...



# Function approximation

How do you design a family of universal models



$$f_{W,b}(x) = W_2 \rho(W_1 x + b_1) + b_2$$

**Theorem (Weierstrass, 1885):** For any continuous function  $f$  on a real interval  $[a, b]$ , for any  $\epsilon > 0$ , there exists a polynomial  $p$  such that  $\|f - p\|_\infty < \epsilon$

universality

**Theorem (Hornik, Stinchcombe, White 1989):** For any continuous function  $f$  on a real hypercube  $[0, 1]^d$ , for any  $\epsilon > 0$ , there exists a two-layer neural net  $p$  such that  $\|f - p\|_\infty < \epsilon$

optimization/GPU-friendly

**Theorem (Telgarsky, 2015):** There exists a family of classification problems indexed by  $k$ , such that to achieve error  $< 1/6$ , two-layer neural nets require  $2^{\Omega(k)}$  nodes while  $2k$ -layer neural nets require  $O(k)$

efficiency

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# Deep learning approaches to learning on graphs

## Non-Euclidean graph space poses unique challenges

- Graph neural networks (GNNs) (Gilmer *et al.*, 2017)
  - Direct parameterization of functions on space of graphs (e.g.  $\mathbb{R}^{n \times n} / S_n$ ) and embellishments (vertex/edge features)
  - Exact invariance (e.g.  $\mathbb{R}^{n \times n} / S_n \rightarrow R$ )

# Graph (convolutional) neural networks

## Deep learning architectures with built-in graph symmetry

- G(C)NN with parameter  $h$ , graph  $G$  and node features  $X$ :

$$\text{GNN}(h, G, X) = X_L(h, G, X)$$

$$X_l(h, G, X) = \rho \left( A_l(h, G) X_{l-1} \right), \quad A_l(h, G) := \sum_{k=0}^K h_{l,k} \text{Adj}(G)^k$$

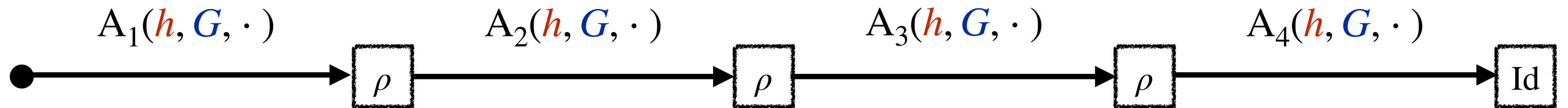
$$X_0(h, G, X) = X$$

- $h$  does not depend on (size of)  $G$



# GNN as operator on node feature

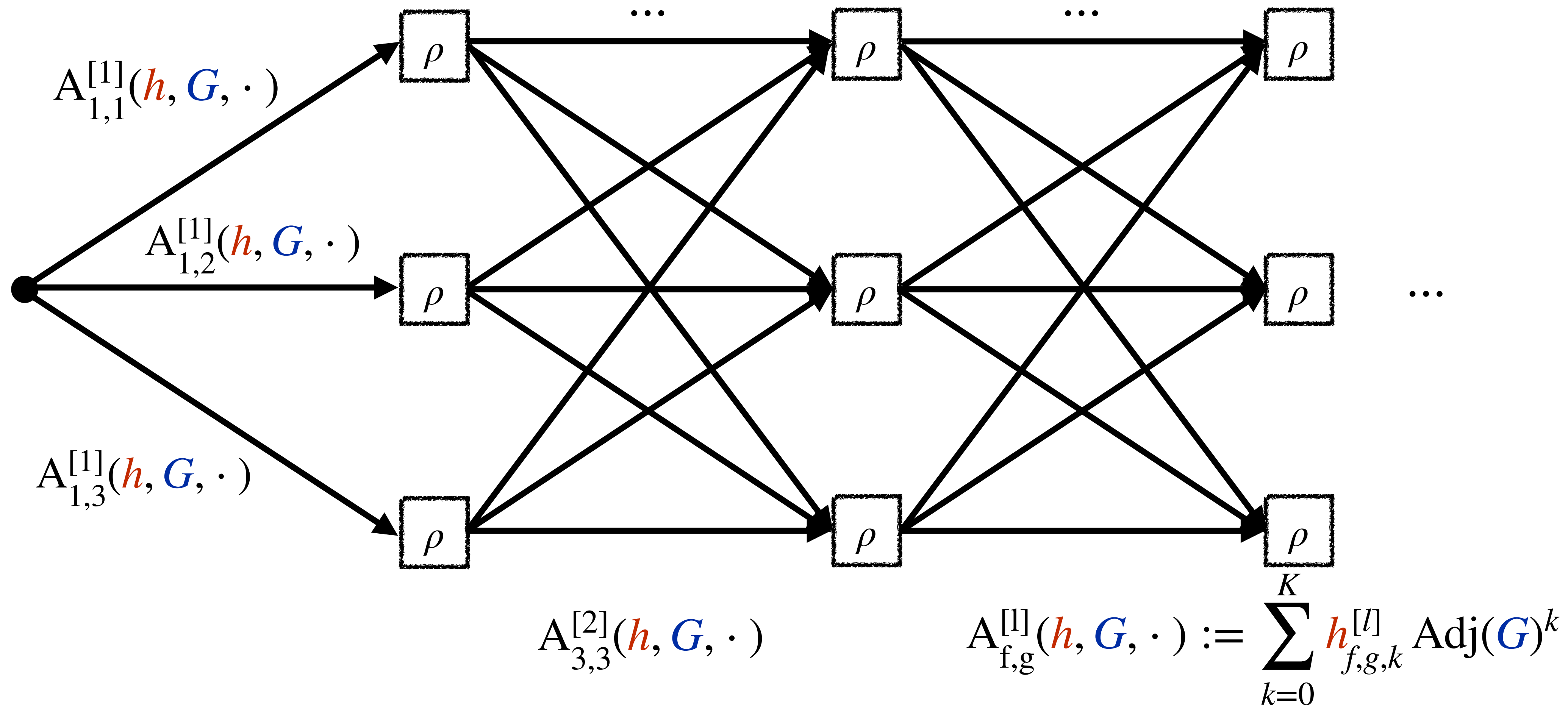
$$\text{GNN}(h, G, \cdot) : \ell^2([n]) \rightarrow \ell^2([n]), \quad n = |V(G)|$$



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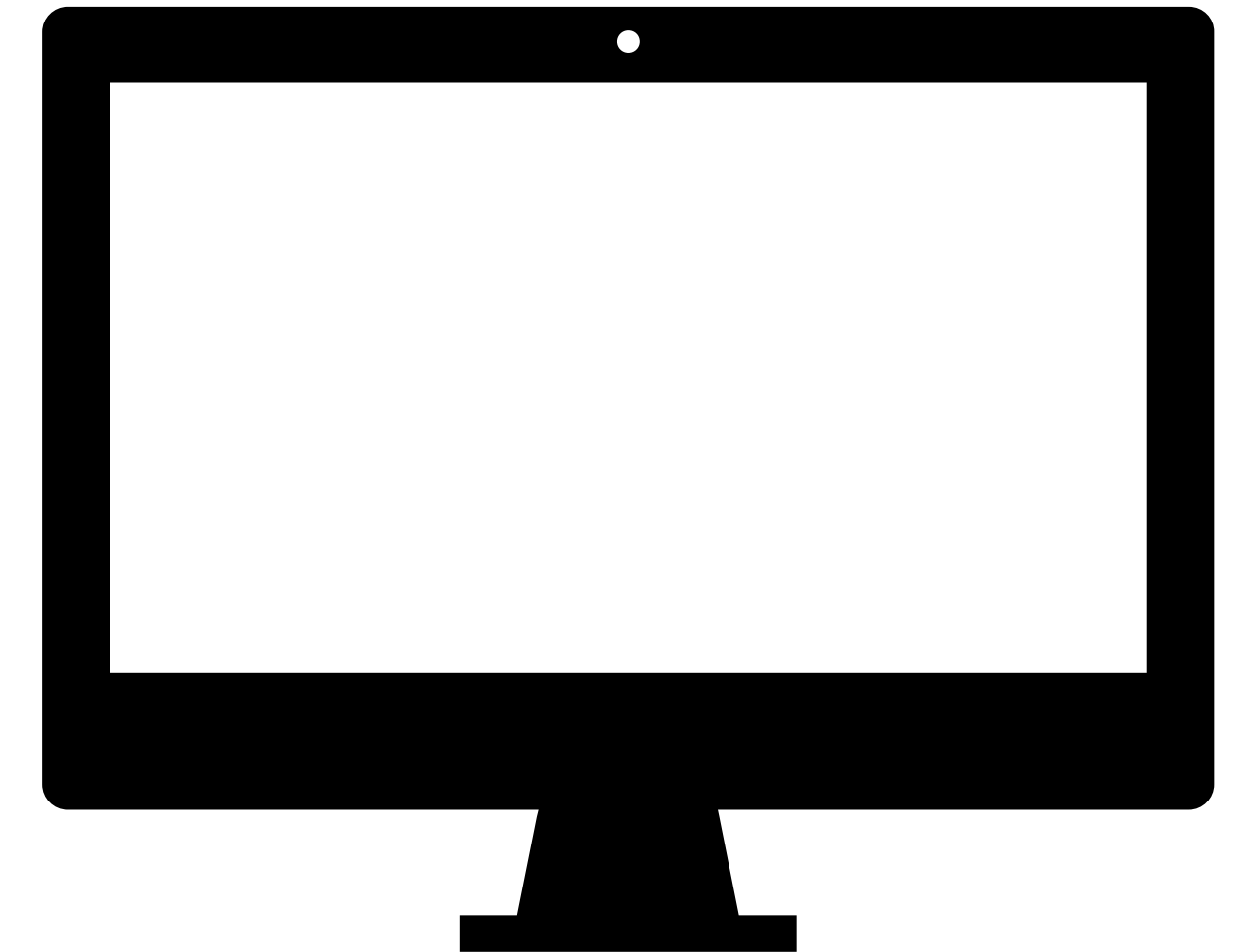


# Is GNN a 'good' model?



Ground truth  
 $f^* : \mathcal{G} \times \mathcal{X} \rightarrow \mathcal{Y}$  that is:  
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...

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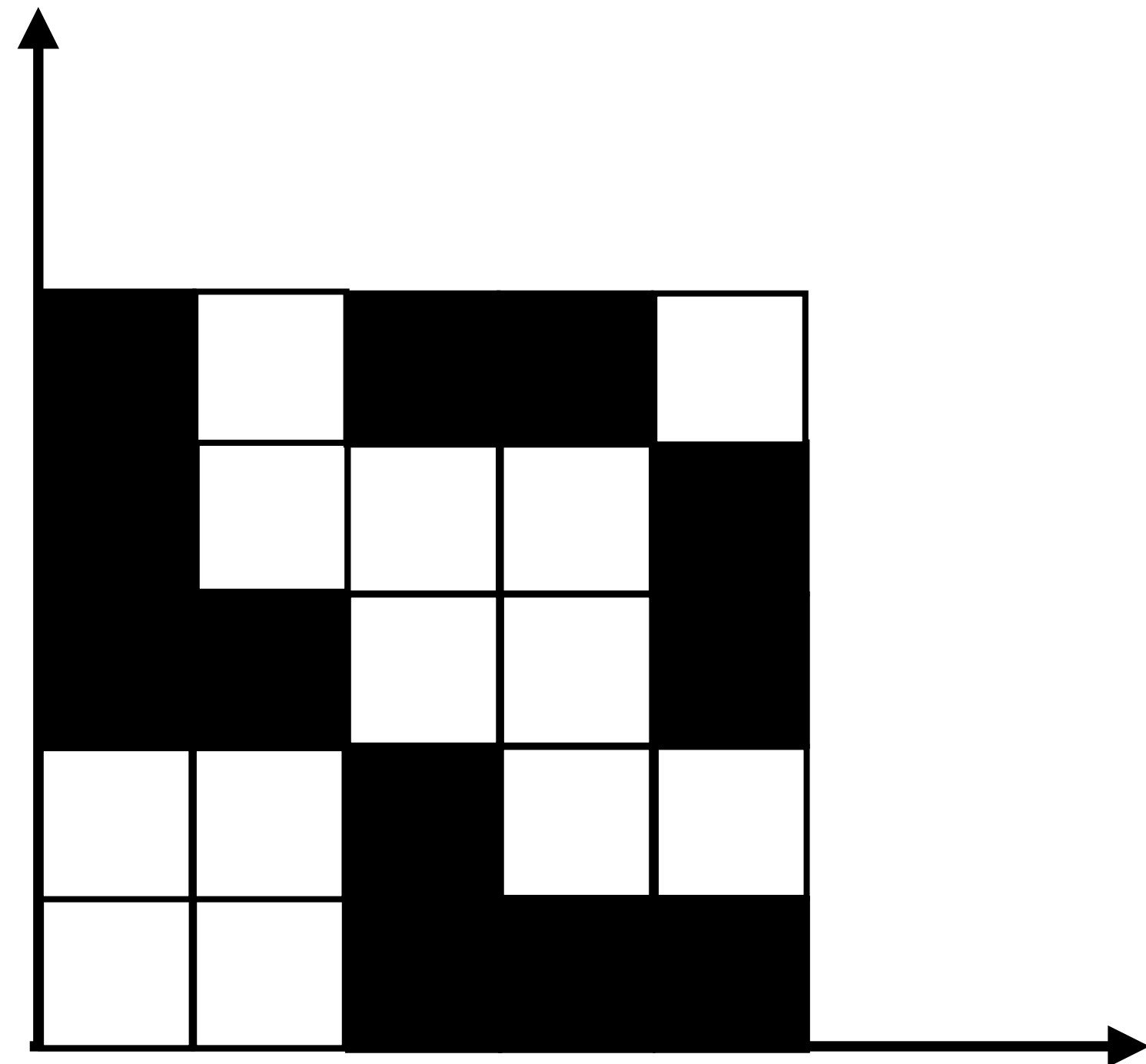
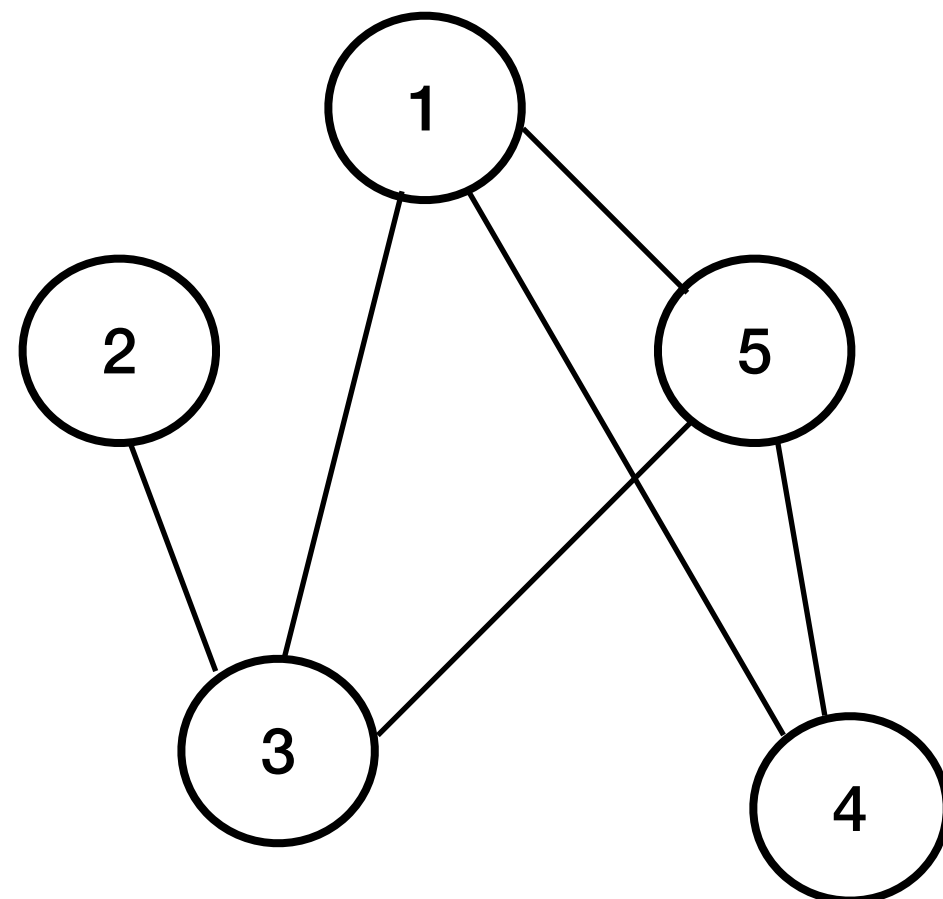
**Universality - Efficiency tradeoff due to hardness of GI**

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# Graphons and the limit of dense graph sequences

- Graphons are symmetric, Lebesgue-measurable functions  
 $W : [0,1] \times [0,1] \rightarrow [0,1]$



# Graphons and the limit of dense graph sequences

Chayes, Borgs, Lovász, Sós, Vesztegombi ~2007

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  2. Finite graph sampler:
    1. Given number of nodes  $n$ , sample  $v_1, \dots, v_n \stackrel{\text{iid}}{\sim} \text{Unif}[0,1]$
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# Most realistic graphs are not dense

## 🌐 Social networks

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Stanford Large Network Dataset Collection

# Approximation for sparse graph sequences

T.L., S. Jegelka (2023)

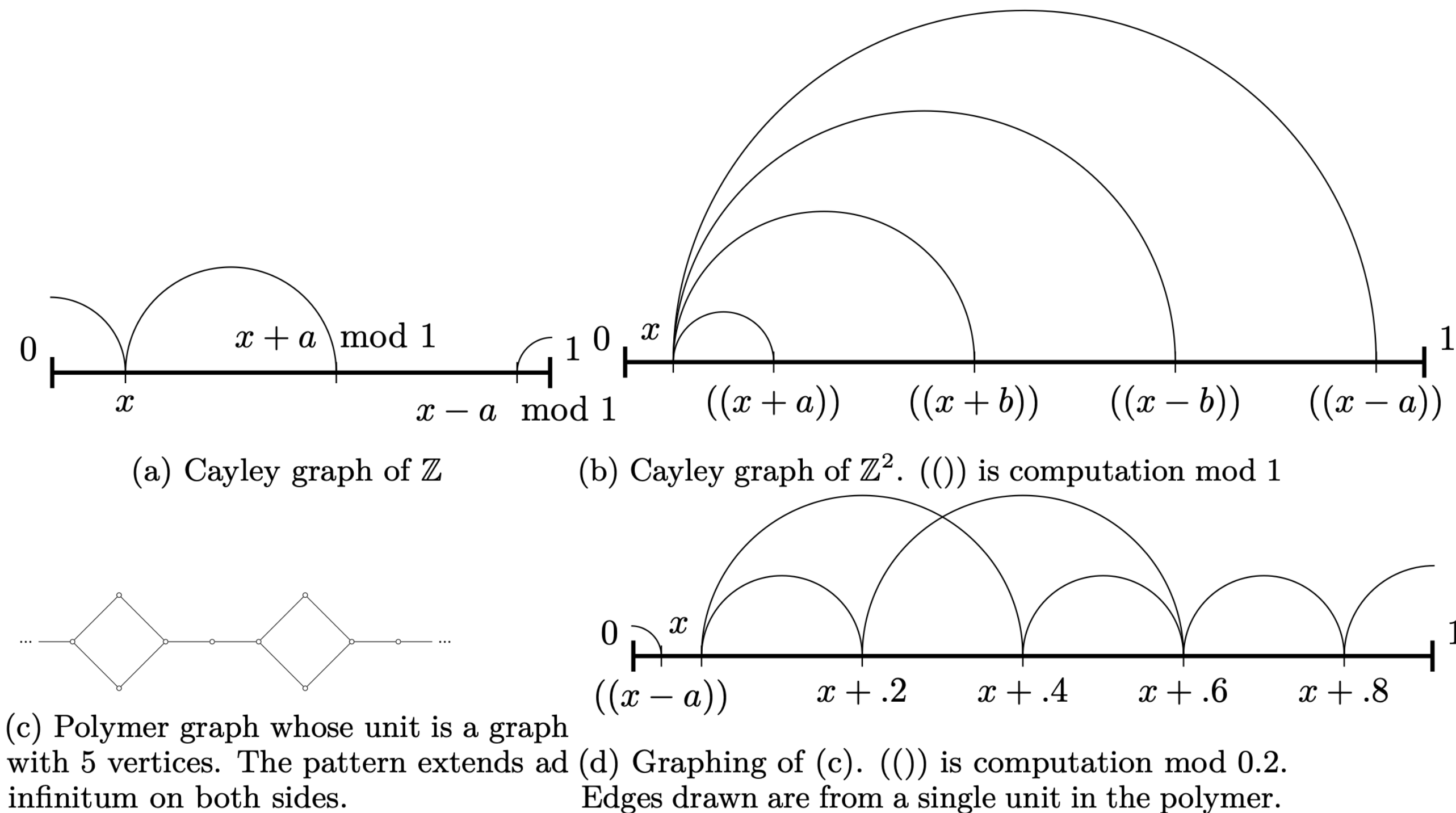
- $d_M$  metric compares **graphop** operators (Backhausz and Szegedy, 2022 + modification to work with continuity assumptions)
- For  $A_n$  a discretization of limit object  $A$

$$d_M(\text{GNN}(h, A_n, \cdot), \text{GNN}(h, A, \cdot)) \leq O\left(n^{-\frac{1}{2}}\right)$$



# Example of sparse graph limit: graphings

Difficult to apply existing spectral techniques.



**Eigengap may not be continuous at the limit!**

Figure 2: Examples of limit objects. The vertex set is the interval  $[0, 1]$ . Example edges are the arcs connecting points on the intervals.  $a$  and  $b$  are distinct irrational numbers. In each graph, edges that miss an endpoint are identified as a single edge connecting the two existing endpoints.

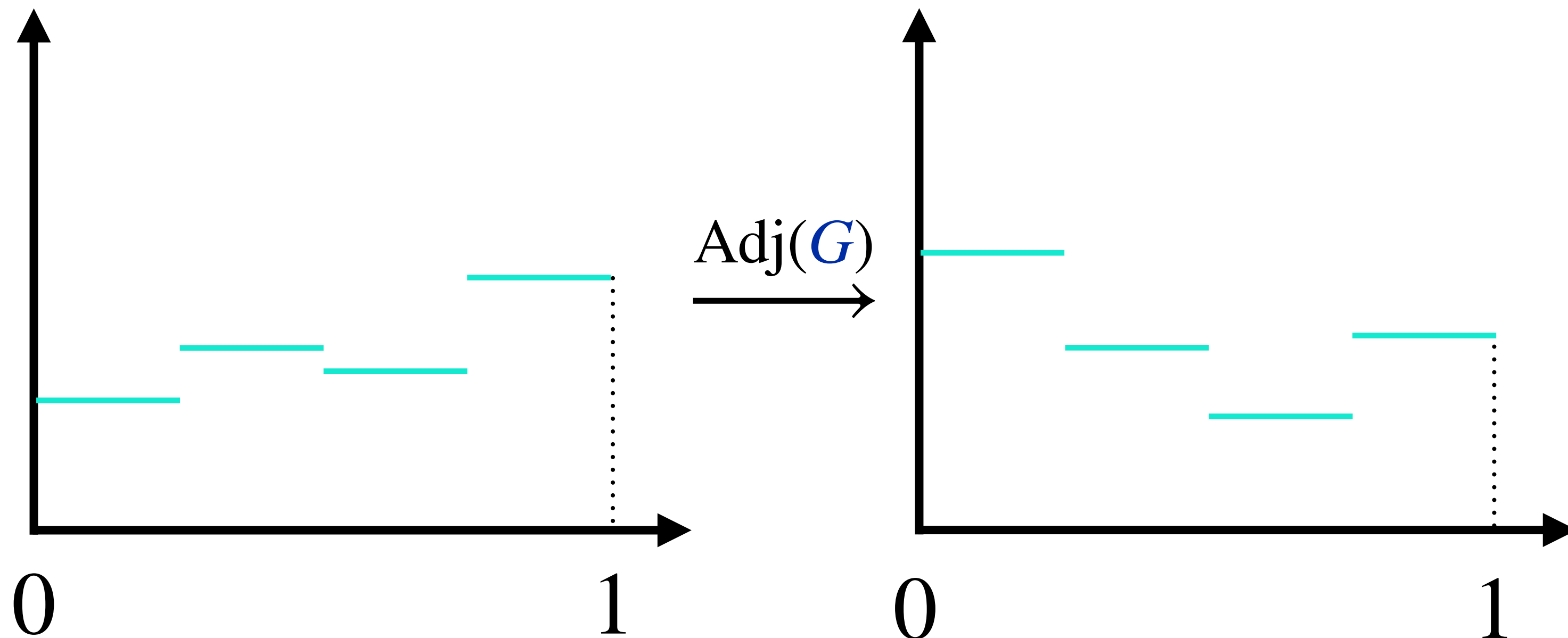
# Assumptions

## Structural assumptions on $\text{Adj}(G)$ instead of spectral

$\text{Adj}(G)$  sends ( $n$ -piece) piece-wise constant function to ( $n$ -piece) piece-wise constant function, for all  $n$  in resolution set  $\mathcal{N}$

OR

$\text{Adj}(G)$  sends **Lipschitz** function to **Lipschitz** function

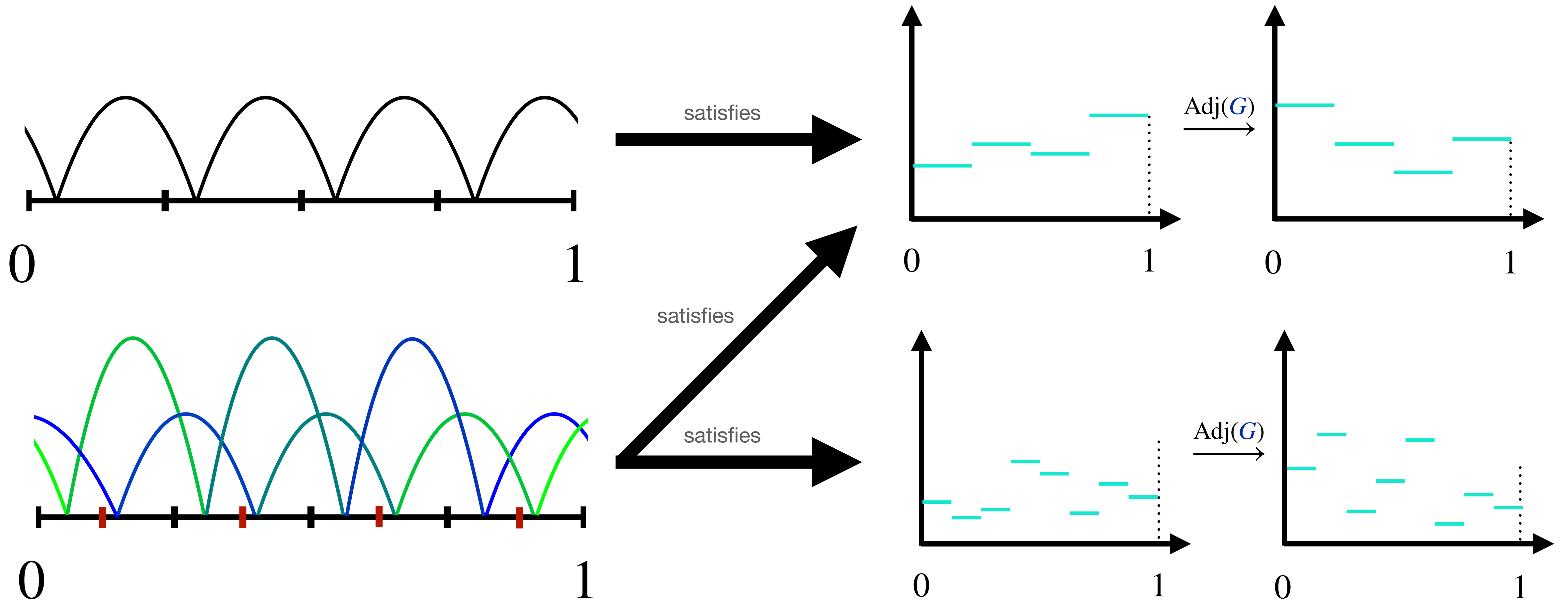


*Other assumptions:*

- $\text{Adj}(G)$  is a Lipschitz operator
- $|h| \leq 1$  point-wise
- $\rho$  is 1-Lipschitz

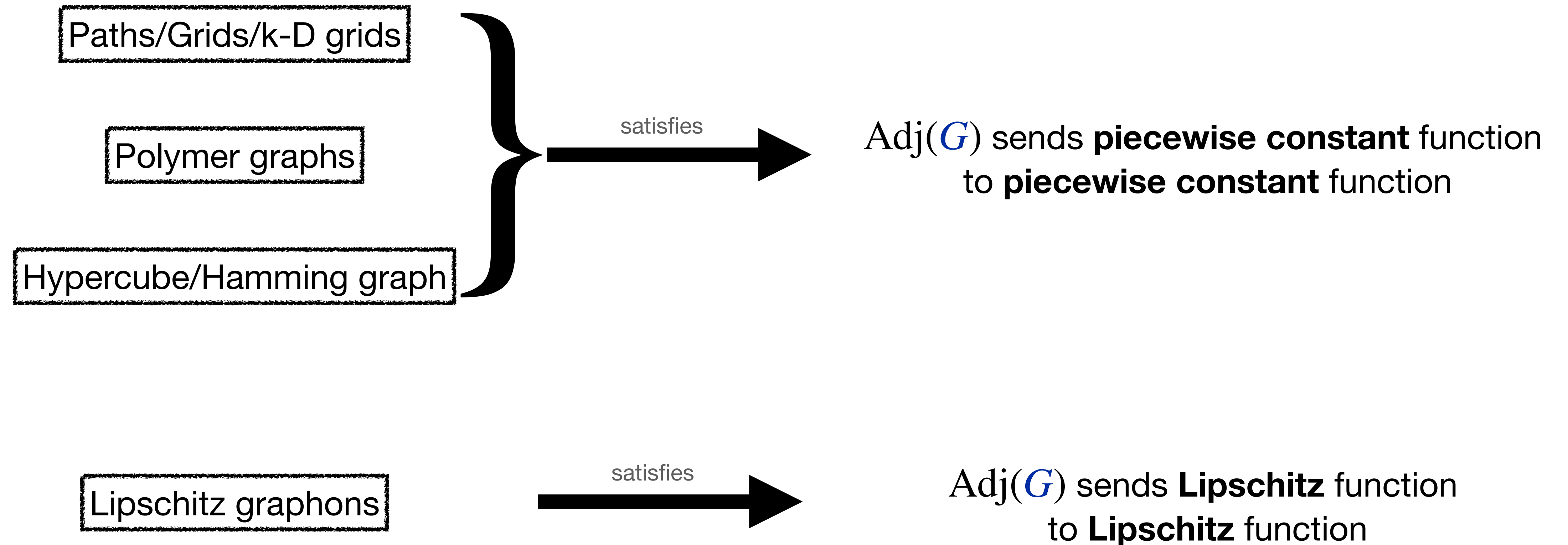
# Example: cycle graphs

$C_4, C_8$  embedded in  $[0,1]$ , resolution set  $\mathcal{N} = \{4,8\}$



# Other examples

## Satisfying structural assumptions



# Approximation for sparse graph sequences

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# Summary

- To obtain approximation theory results, additional structures are needed on the space of graphs
- We prove an approximation result for GNNs by graph limit
- Unlike dense graphs, sparse graph limits can be pathological, which was circumvented by enforcing structural assumptions

**Thank you**

**Q&A**





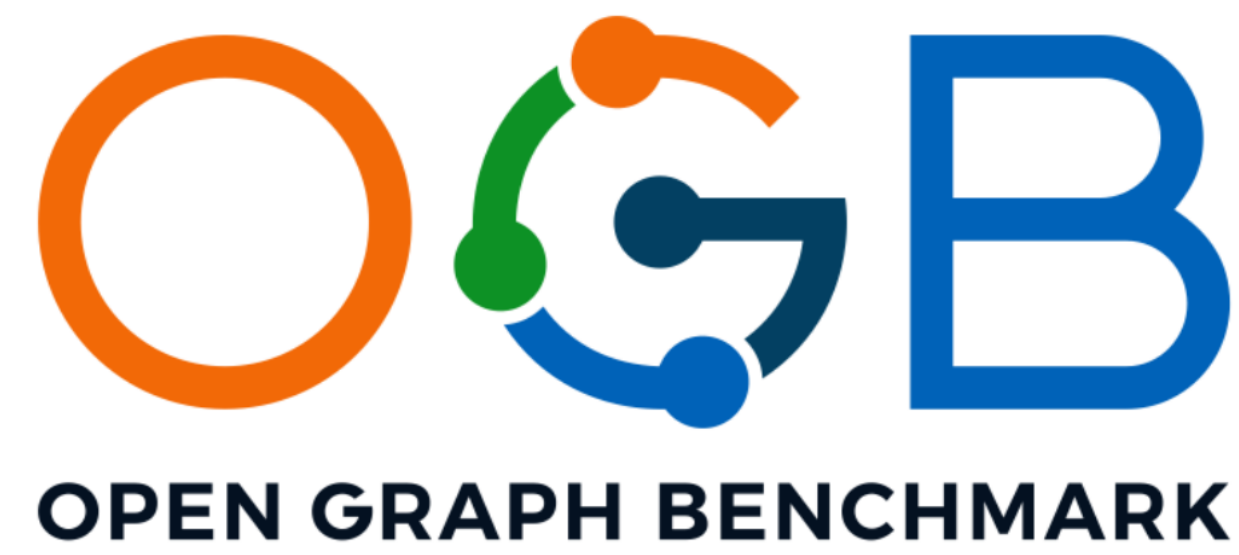
# Machine learning on graph structures

- Social networks
  - Community detection
  - Link prediction



Stanford Large Network  
Dataset Collection

- Molecular graphs
  - Property prediction
  - Geometry prediction



- Traditional tasks in vision (pixel graphs) or NLP (path graphs)

# Deep learning approaches to graph learning

- Graph neural networks (GNNs) (Gilmer *et al.*, 2017)
  - Direct parameterization of functions on space of graphs (e.g.  $\mathbb{R}^{n \times n} / S_n$ ) and embellishments (vertex/edge features)
  - Exact invariant (e.g.  $\mathbb{R}^{n \times n} / S_n \rightarrow R$ ) or equivariant (e.g.  $\mathbb{R}^{n \times n} / S_n \rightarrow R^n / S_n$ )
- Graph transformers (Veličković *et al.*, 2018)
  - Graph information added to input of transformer (self attention layers)
- Mixed architectures

# Graph (convolutional) neural networks

## Deep learning architectures with built-in graph symmetry

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$$\text{GNN}(h, G, X) = X_L(h, G, X)$$

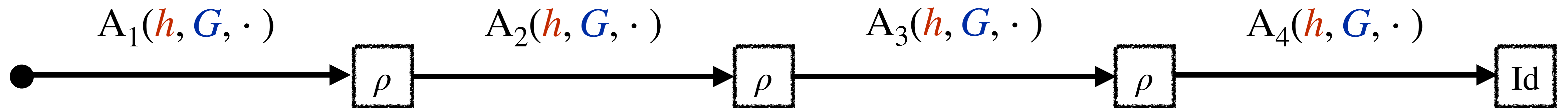
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# GNN as operator on node feature

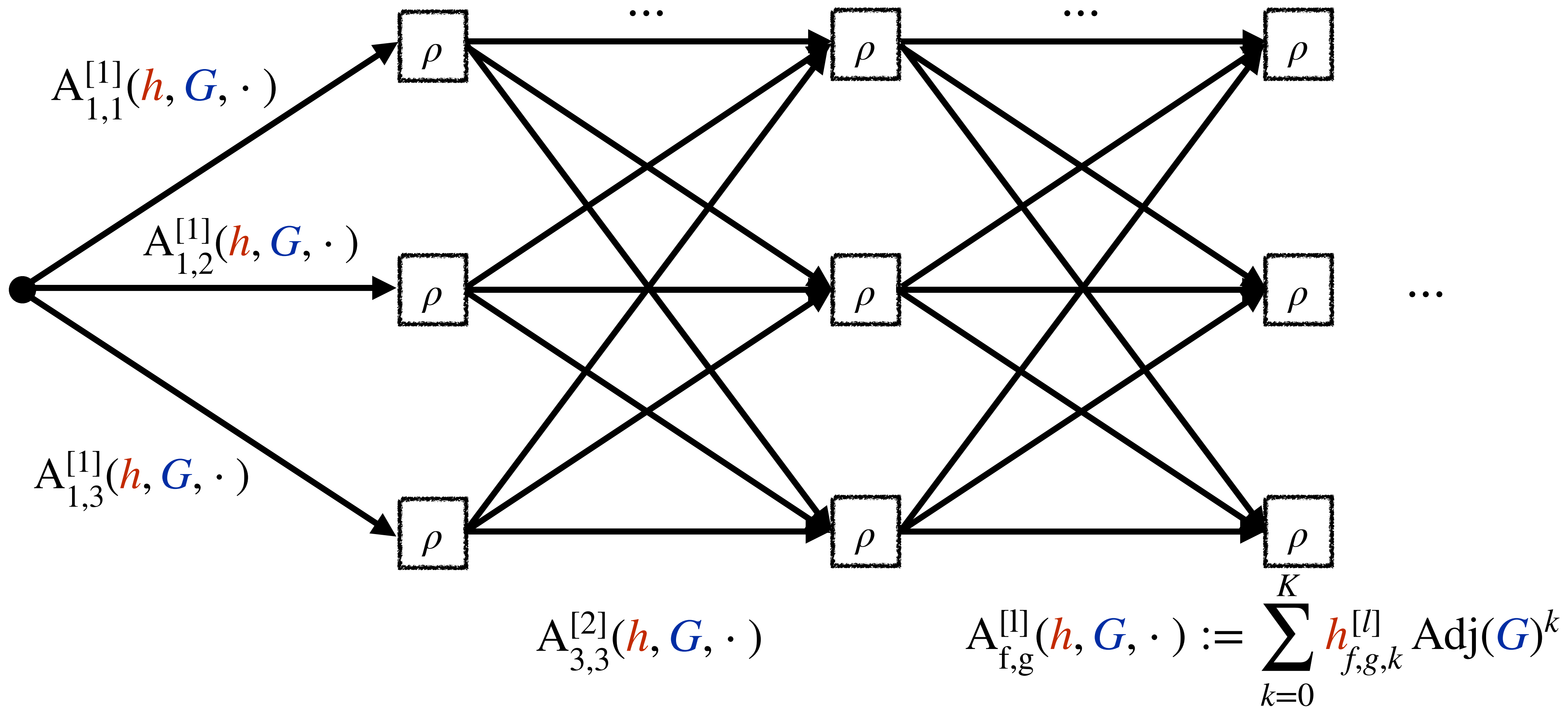
$$\text{GNN}(h, G, \cdot) : \ell^2([n]) \rightarrow \ell^2([n]), \quad n = |V(G)|$$



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$$\text{GNN}(h, G, \cdot) : \ell^2([n]) \rightarrow \ell^2([n]), \quad n = |G| := |V(G)|$$



# Transferability of GNNs

**GNNs on ‘similar’ graphs are also ‘similar’**

- Train on small graphs and test on large graphs
- Motivations:
  - Explanation for pretraining performance
  - Training vs evaluation compute difference
  - Out-of-distribution generalization/extrapolation

# Transferability of GNNs

GNNs on 'similar' graphs are also 'similar'

$$d_{\text{GNN}}(\text{GNN}(h, G_1, \cdot), \text{GNN}(h, G_2, \cdot)) \quad \text{vs} \quad d_G(G_1, G_2)$$

# Size transferability of GNNs

GNNs on 'similar' graphs of **different sizes** are also 'similar'

Sequence of graphs  $F_1, F_2, F_3, F_4, \dots$  of increasing size  $n_1, n_2, n_3, n_4, \dots$

$d_{\text{GNN}}(\text{GNN}(h, G_1, \cdot), \text{GNN}(h, G_2, \cdot))$  vs decreasing\_fn( $|G_1|, |G_2|$ )

e.g.  $1/|G_1|^c + 1/|G_2|^c, c > 0$



# Size transferability of GNNs

A limit approach ( $F_n \xrightarrow{n \rightarrow \infty} G^\infty$  in some limit)

$$d_{\text{GNN}}(\text{GNN}(h, G_1, \cdot), \text{GNN}(h, G_2, \cdot))$$

$$\leq$$

$$d_{\text{GNN}}(\text{GNN}(h, G_1, \cdot), \text{GNN}(h, G^\infty, \cdot)) \quad \text{vs} \quad O(1/|G_1|^c)$$

$$+$$

$$d_{\text{GNN}}(\text{GNN}(h, G_2, \cdot), \text{GNN}(h, G^\infty, \cdot)) \quad \text{vs} \quad O(1/|G_2|^c)$$

approximation theorems

# This paper

We show **size** transferability for different graph sequences, in particular sparse graphs

Via approximation bound:  $d_{\text{GNN}}(\text{GNN}(h, G_1, \cdot), \text{GNN}(h, G^\infty, \cdot)) \leq O(1/|G_1|^c)$

	Sparse		Dense
	Bounded-degree	Relatively-sparse	
Number of edges	$\Theta(n)$	$\Theta(n \log n)$	$\Theta(n^2)$
Examples covered under our assumptions	infinite grids, polymer graphs	hypercubes, Hamming graphs	graphons
Graphons (Ruiz et al., 2023a)			$O(n^{-1})$
Unbounded graphons (Maskey et al., 2023)			inexplicit
Random graph model (Keriven et al., 2020)		$O((\log n)^{-1/2})$	$O(n^{-1/2})$
Spectral methods (1 layer) (Levie et al., 2022)		inexplicit	inexplicit
Graphings (1 layer) (Roddenberry et al., 2022)	inexplicit		

Table 1: Summary of our results compared to related work. Quantitative results (e.g.  $O(n^{-1})$ ) upper-bound the distance between GNNs on sampled graphs of size  $n$  and the limiting object in term of  $n$  (in an appropriate metric and limit notion). Empty cells are graph models where the approaches in the corresponding papers do not apply to or give trivial bounds (e.g. bounds that compare to a constant-0 graphon). "Inexplicit" refers to asymptotic results where rates of convergence is not explicit.

# Graph limit: graphons

Chayes, Borgs, Lovász, Sós, Vesztegombi ~2007

- Graphons are symmetric, Lebesgue-measurable functions  
 $\mathbf{W} : [0,1] \times [0,1] \rightarrow [0,1]$
- Interpretation:
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# Transferability via graphons

Ruiz, Chamon, Ribeiro, 2023

$$\| \text{GNN}(h, G_1, \cdot) - \text{GNN}(h, G_2, \cdot) \|_{L_2} \leq O(|G_1|^{-1} + |G_2|^{-1}) + \epsilon$$

Note: It is possible to optimize Ruiz *et al.*'s bound to get rid of the  $\epsilon$  but incur a slower rate

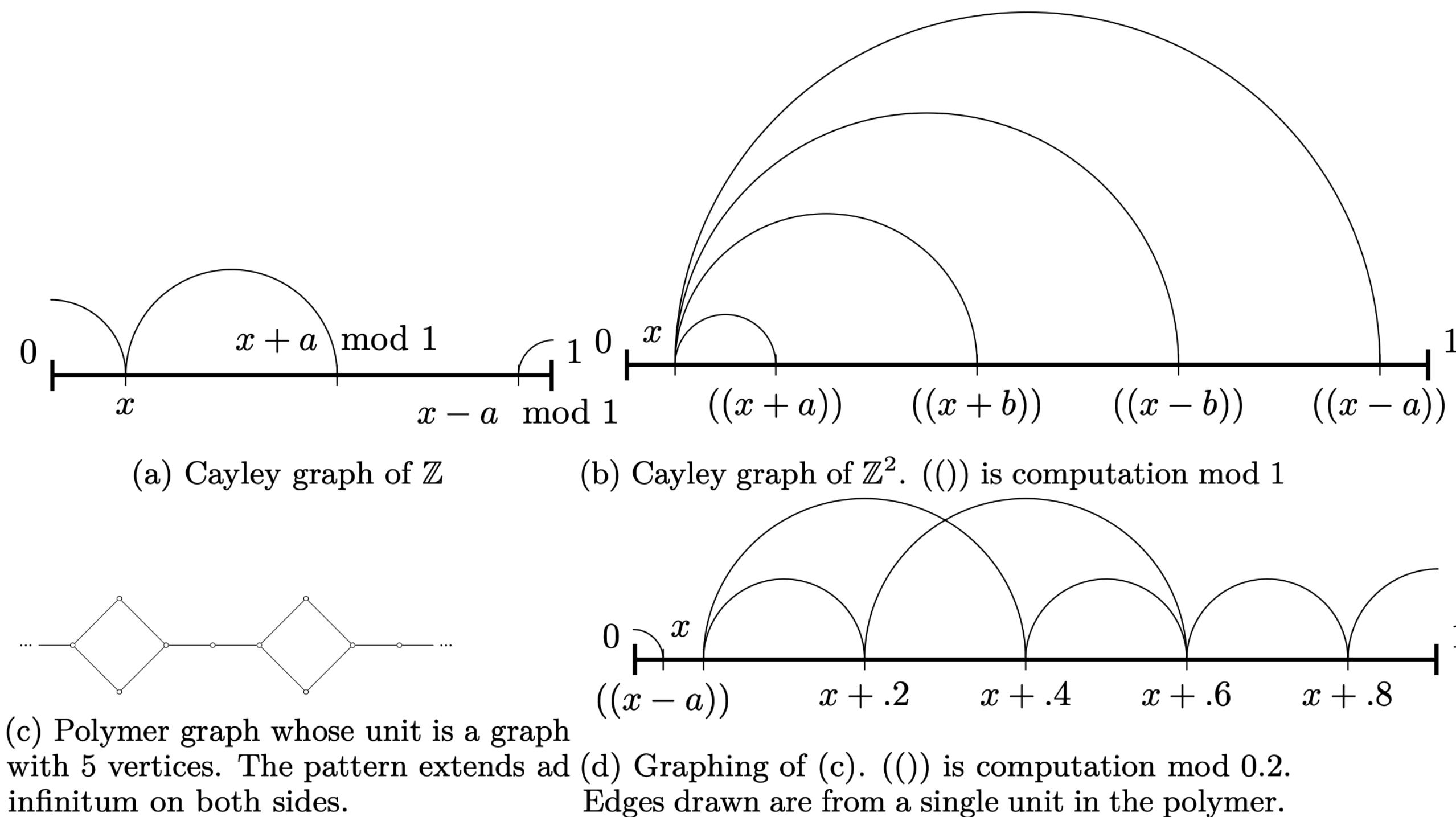
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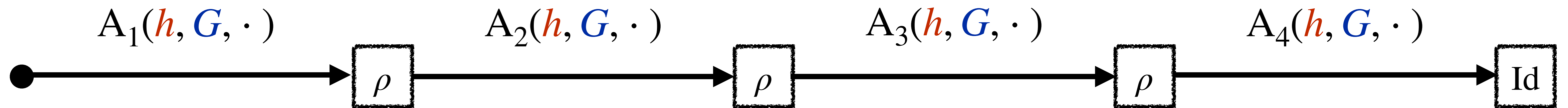


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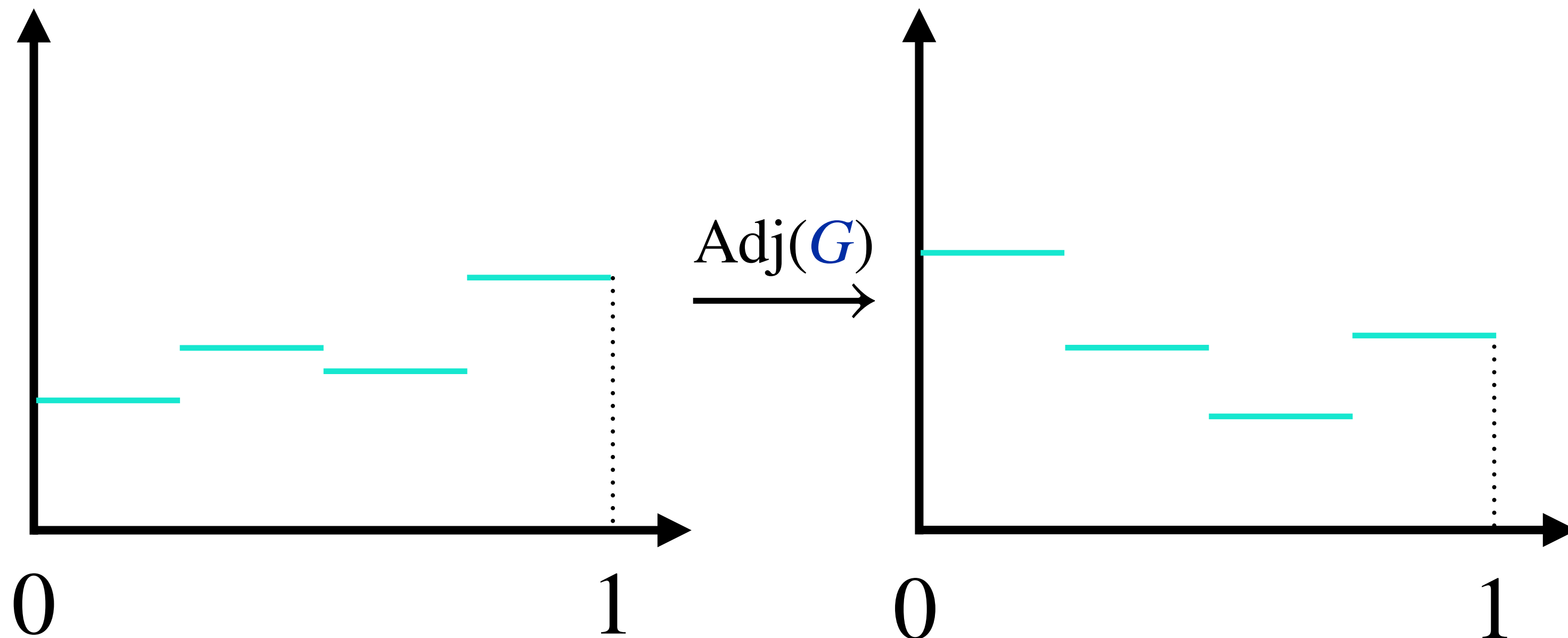
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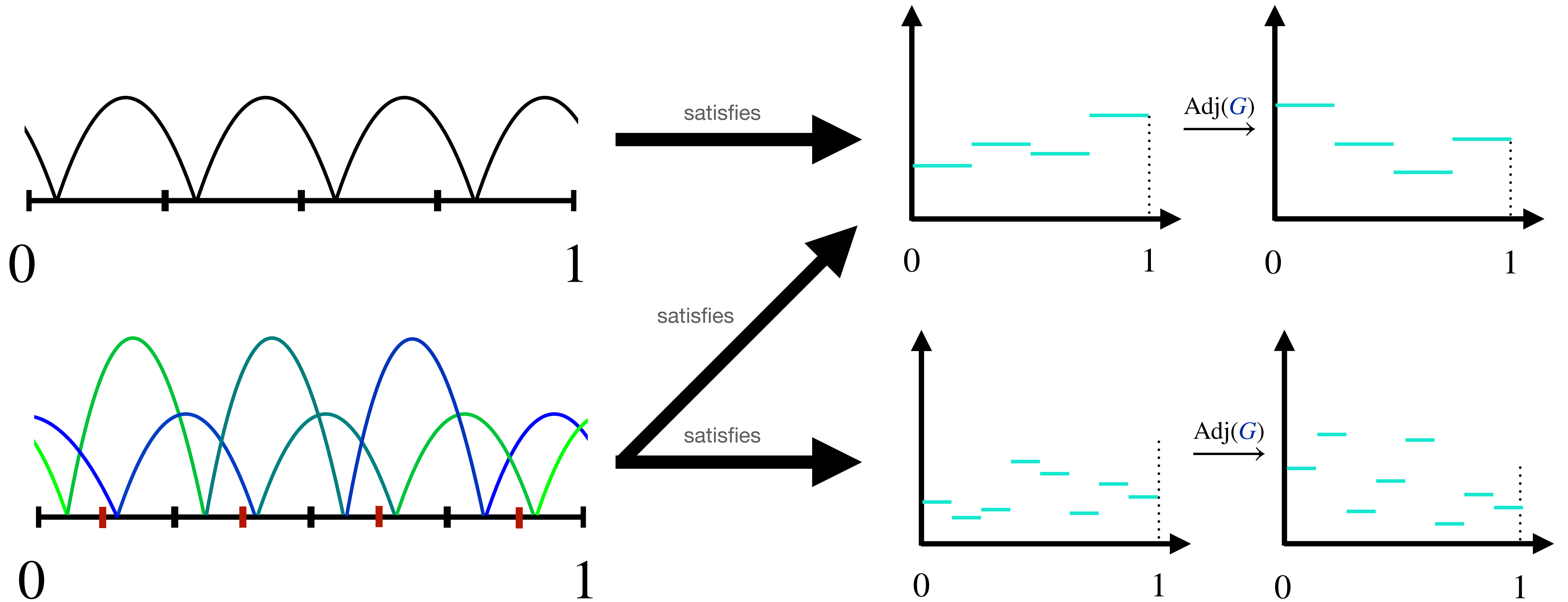
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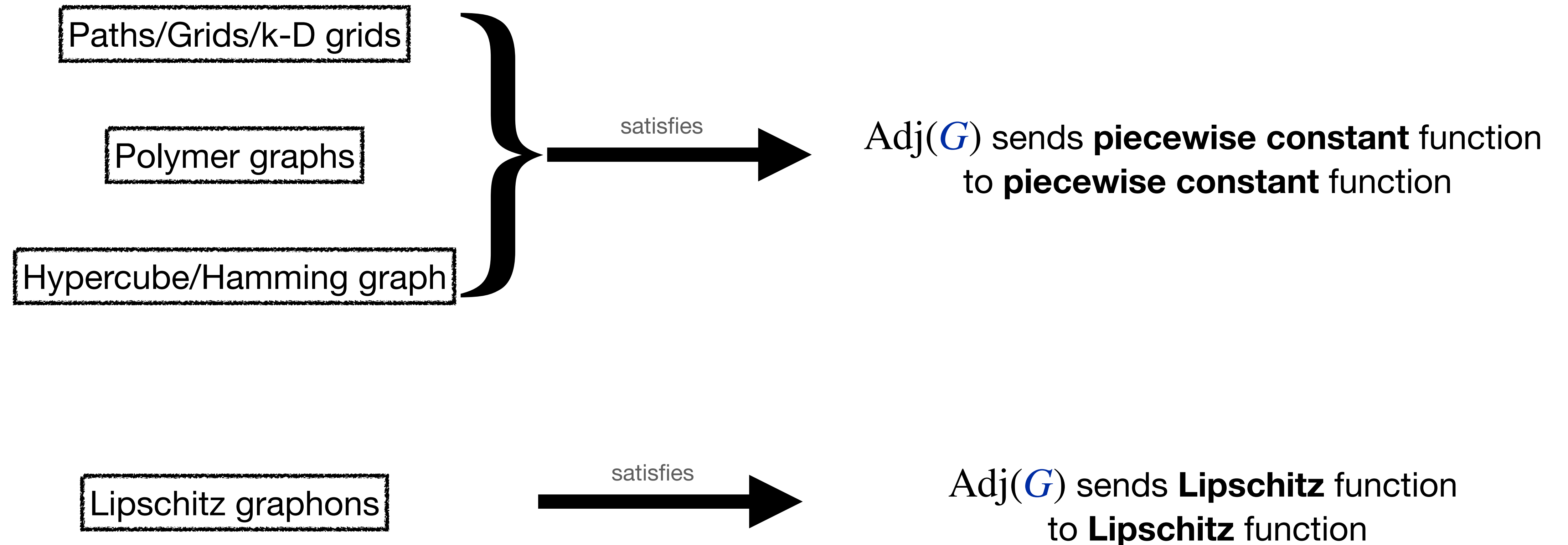
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# Other examples

## Satisfying structural assumptions



# Main theorem: Approximation and Size transferability via graphop (Backhausz and Szegeedy 2022: action convergence)

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- For  $A_n$  a discretization of limit object  $A$

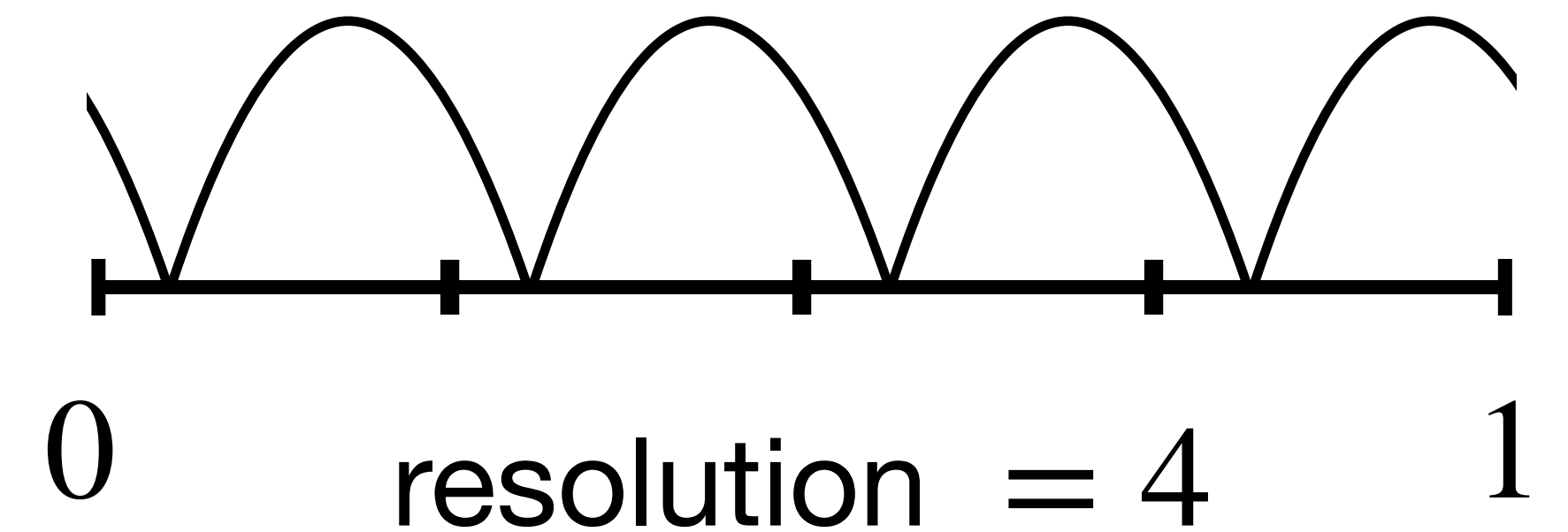
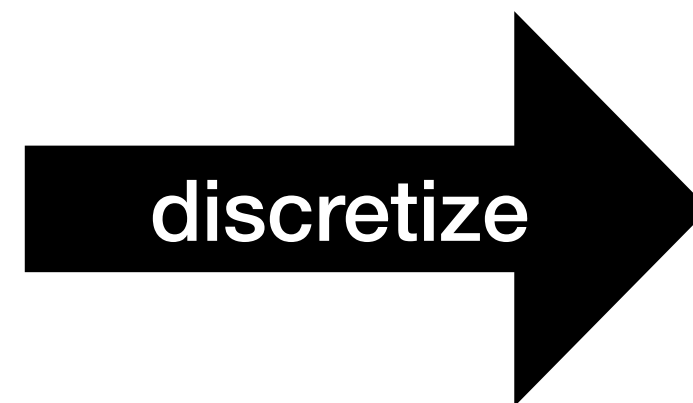
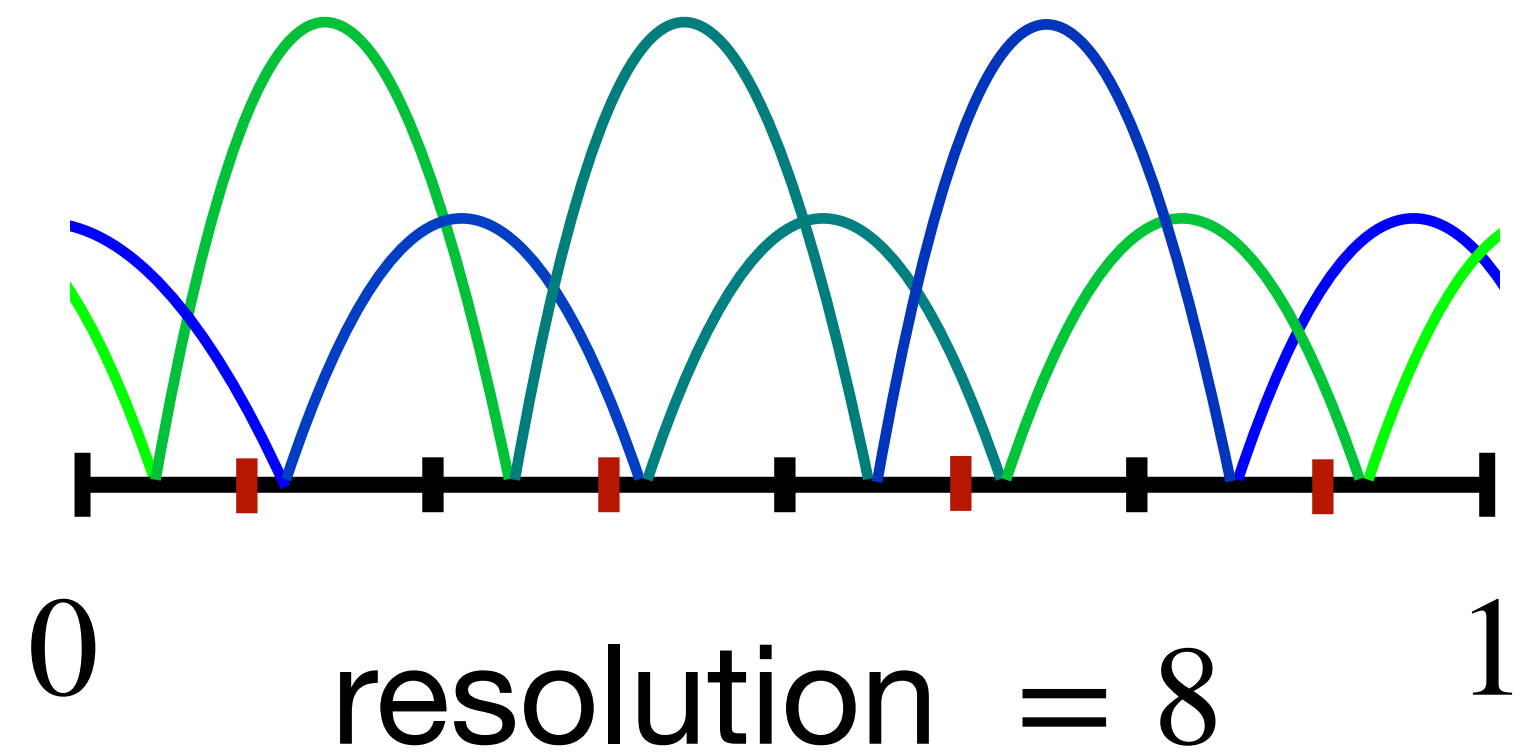
$$d_M(\text{GNN}(h, A_n, \cdot), \text{GNN}(h, A, \cdot)) \leq O\left(n^{-\frac{1}{2}}\right)$$

- For  $A_m, A_n$  two different discretizations of the same limit object of size  $m, n$

$$d_M(\text{GNN}(h, A_m, \cdot), \text{GNN}(h, A_n, \cdot)) \leq O\left(m^{-\frac{1}{2}} + n^{-\frac{1}{2}}\right)$$

# Discretizing adjacency operator/GNNs

‘Averaging connections’



More generally,

$$A_m X(v) := m \int_{v-\frac{1}{m}}^v (A\tilde{X})d\lambda, \text{ for all } v \in [m]/m, \quad \text{GNN}_m(h, A, X) := \text{GNN}(h, A_m, X),$$

# Summary

- We prove an approximation and size transferability result for GNNs by graph limit.
- Unlike dense graphs, sparse graph limits can be pathological.
- By enforcing structural assumptions, our result works for sparse and dense graph limits.

# Future directions

- Relaxing assumptions
  - Warning: unconditional approximation theorem solves group theoretic open questions (e.g. existence of non-sofic groups) - Backhausz and Szegedy.
  - But other than that?
- Graph transformer
- Optimization of sequence transformer via Kuramoto model

# Q & A

- Thank you for your attention.

**Back up slides**



# Optimization of transformers via Kuramoto model

## A MATHEMATICAL PERSPECTIVE ON TRANSFORMERS

BORJAN GESHKOVSKI, CYRIL LETROUT, YURY POLYANSKIY,  
AND PHILIPPE RIGOLLET

ABSTRACT. Transformers play a central role in the inner workings of large language models. We develop a mathematical framework for analyzing Transformers based on their interpretation as interacting particle systems, with a particular emphasis on long-time clustering behavior. Our study explores the underlying theory and offers new perspectives for mathematicians as well as computer scientists.

*Part 3: Further questions.* We propose potential avenues for future research, largely in the form of open questions substantiated by numerical observations. We first focus on the case  $d = 2$  (Section 7) and elicit a link to Kuramoto oscillators. We

## Graphop Mean-Field Limits for Kuramoto-Type Models

Marios-Antonios Gkogkas\* and Christian Kuehn \*

December 16, 2020

For dense graphs converging to graphon limits, one also knows that mean-field approximation holds for certain classes of models, e.g., for the Kuramoto model on graphs. Yet, the space of intermediate density and sparse graphs is clearly extremely relevant. Here we prove that the Kuramoto model can be approximated in the mean-field limit by far more general graph limits than graphons.

# Graphings and Benjamini-Schramm convergence

## Local convergence of bounded-degree graphs

- $(X, \mathcal{B}, \nu)$  Borel probability measure on topological space  $X$ . Graphings are graphs with vertex set  $X$  and Borel edge set  $E \subset X \times X$  with a symmetric constraint.
- Rooted distance:  $d(G_1, G_2) = 1/k$  where  $k$  is the smallest number such that the  $k$ -neighborhood of  $G_1$  and  $G_2$  around their roots are isomorphic.
  - Theorem: space of rooted graph with rooted distance is compact
  - Cor: space of Borel probability measures on rooted graphs is compact in weak topology
  - BS convergence: embed graphs as Borel measures that is uniform over root, for each radius

# Graphop

Backhausz and Szegedy, 2022

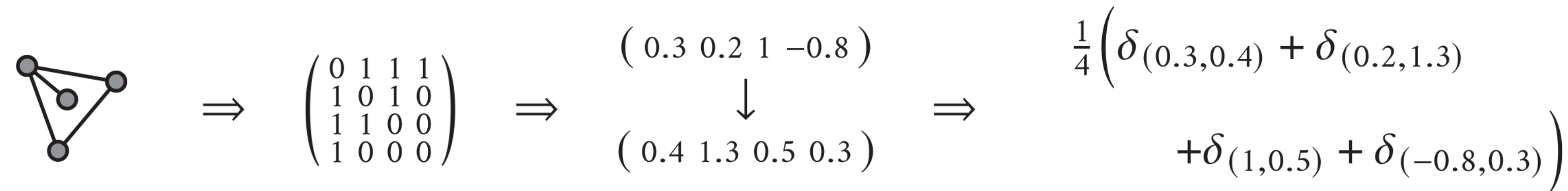


Figure: Backhausz and Szegedy

Figure 2: Graph  $\Rightarrow$  operator  $\Rightarrow$  action  $\Rightarrow$  measure (computing an element in the 1-profile of a graph).

- P-operators: linear bounded (in operator norm) operators

- k-profile: the set  $\mathcal{S}_k(A)$  of all possible probability measures of the form  $\mathcal{D}(v_1, \dots, v_k, Av_1, \dots, Av_k) = \frac{1}{n} \sum_{j=1}^n \delta_{(v_{1,j}, \dots, v_{k,j}, [Av_1]_j, \dots, [Av_k]_j)}$

- Hausdorff distance between closed sets of probability measures:  $d_H(X, Y) = \max(\sup_x \inf_y d(x, y), \sup_y \inf_x d(x, y))$

- Action convergence metric:  $d_M(A, B) = \sum_{k=1}^{\infty} 2^{-k} d_H(\mathcal{S}_k(A), \mathcal{S}_k(B))$

# Small experiment for rate

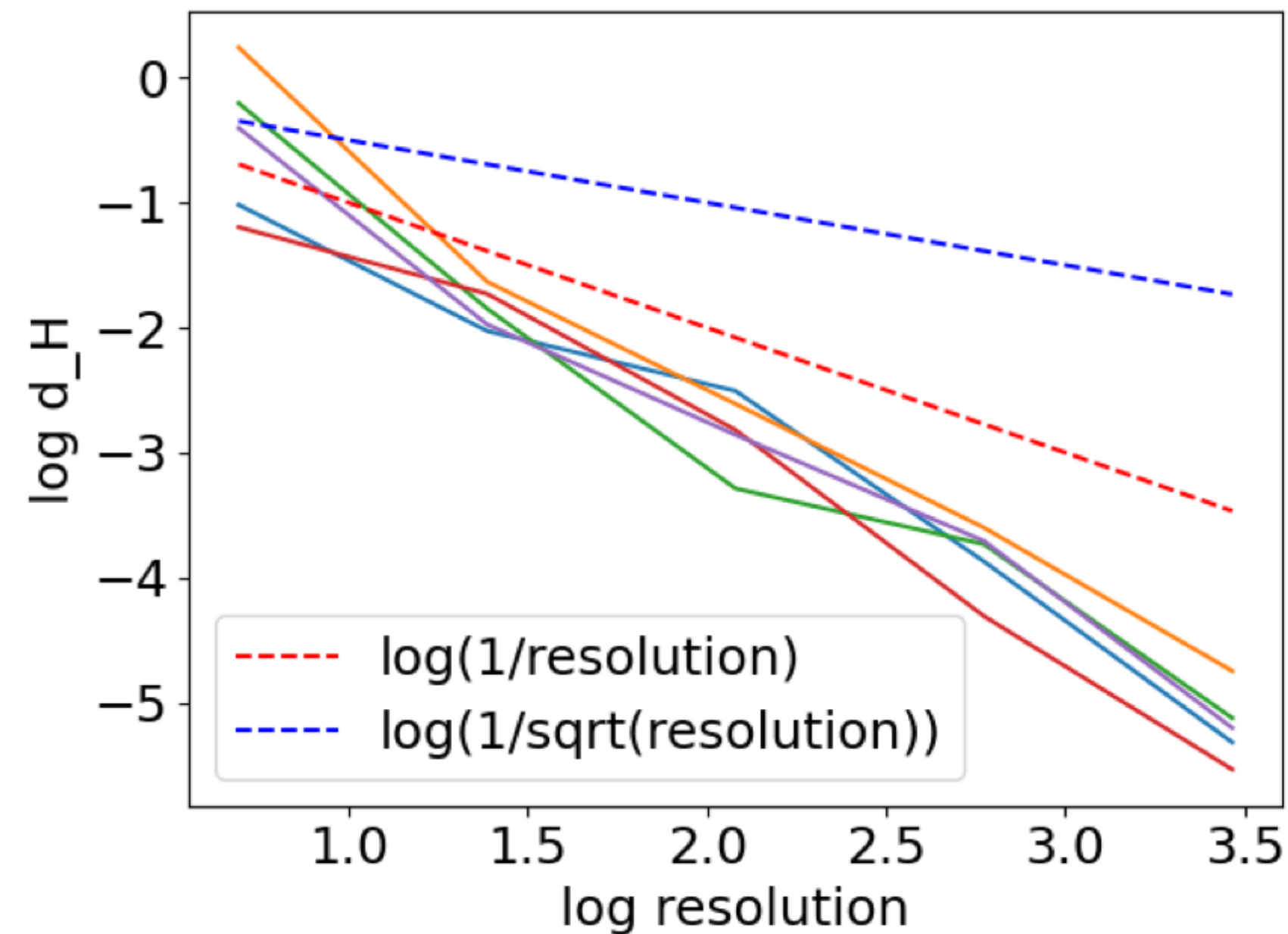


Figure 1: Hausdorff metric between samples from 1-profiles of 2-hidden-layer GNN on finite polymer graphs vs on large polymer graphs (see Appendix A for polymer graphs). The GNN uses GSO  $A_n^2 + A_n$  where  $A_n$  is the normalized adjacency matrix on  $n$  nodes and ReLU nonlinearities at each layer. Different solid lines are different random draws of functions that make up the estimated 1-profile. See Appendix A for details.

# Graphop neural network as P-operator

- There is no requirement for P-operators to be linear!
- Action convergence can be defined for nonlinear operators

**Conjecture 1** (Action convergence of graphop neural networks). *Let  $(A_n)_{n \in \mathbb{N}}$  be an action convergent sequence of graphops. Then  $(\Phi(h, A_n, \cdot))_{n \in \mathbb{N}}$  is an action convergent sequence of P-operators.*